

# NIM

MATH CIRCLE (ADVANCED) 12/02/2012

Recall from last time the idea of dividing a game into **P-positions** and **N-positions**.

A P-position is one that is winning for the previous player (that is, the one who just moved) while a N-position is one that is winning for the next player.

Therefore, if a game starts in a P-position, Player II wins, while if the game starts in a N-position, Player I wins.

The game of Nim: Suppose we have  $k$  piles of stones, with  $n_1$  stones in the first pile,  $n_2$  stones in the second pile, etc. Player I and Player II take turns removing any number (at least one!) stone from a single pile. The player who takes the last stone wins.

We will write the above game as  $(n_1, n_2, \dots, n_k)$ -Nim.

0) a) List the P-positions and N-positions for Nim played with only one pile ( $k = 1$ ).

The only P-position is  $(0)$ .

b) For what values of  $n_1$  does Player I win?

Player I wins all games.

1) a) Play Nim with piles of size 6 and 4 (i.e.  $(6, 4)$ -Nim). Who wins?

Player I wins.

b) Explain the winning strategy.

Player I first moves to position  $(4, 4)$ . Then he plays so that both piles always have the same number after his turn.

2) a) Characterize the P-positions and N-positions for  $(n_1, n_2)$ -Nim.

$(n_1, n_2)$  is a P-position if and only if  $n_1 = n_2$ .

b) For what pairs  $(n_1, n_2)$  does Player I win?

Player I wins  $(n_1, n_2)$ -Nim if and only if  $n_1 \neq n_2$ .

The Nim-sum (denoted by  $\oplus$ ) of two numbers is computed as follows. Convert the numbers to binary, then add (in binary) WITHOUT carrying any digits, and finally convert back to decimal form. For example,

$$60 \oplus 47 \rightarrow \begin{array}{r} 1 \ 1 \ 1 \ 1 \ 0 \ 0 \\ 1 \ 0 \ 1 \ 1 \ 1 \ 1 \\ \hline 0 \ 1 \ 0 \ 0 \ 1 \ 1 \end{array} \rightarrow 10011_2 = 19.$$

3) Compute the following Nim-sums:

a)  $6 \oplus 4$

2

b)  $12 \oplus 12$ 

0

c)  $25 \oplus 8$ 

17

d)  $3 \oplus 4 \oplus 5$ 

2

e)  $19 \oplus 60 \oplus 47$ 

0

f)  $2 \oplus 2 \oplus 2 \oplus 2$ 

0

g)  $5 \oplus 2 \oplus 2 \oplus 2$ 

7

4) Write out P-positions and N-positions for  $(n_1, n_2, n_3)$ -Nim for  $n_1 + n_2 + n_3 \leq 6$ . Think of ways to minimize the number of positions you need to write out!

First note that without loss of generality  $n_1 \geq n_2 \geq n_3$  (why?). Further, if  $n_3 = 0$  see 2)a). Therefore we also assume  $n_3 > 0$ .

The only P-position satisfying the above are:  $(3, 2, 1)$ .

The N-positions are  $(1, 1, 1)$ ,  $(2, 1, 1)$ ,  $(2, 2, 1)$ ,  $(3, 1, 1)$ .

5) a) Prove that  $a \oplus b = 0$  if and only if  $a = b$ .

If  $a = b$  all the binary digits are the same, so  $a \oplus b = 0$ . If  $a \neq b$  then (at least) one pair of binary digits do not match, hence  $a \oplus b \neq 0$ .

b) Restate 2) in terms of Nim-sums.

$(n_1, n_2)$  is a P-position if and only if  $n_1 \oplus n_2 = 0$ .

6) a) Characterize the P-positions and N-positions for  $(n_1, n_2, n_3)$ -Nim.

See 8).

b) For what triples  $(n_1, n_2, n_3)$  does Player I win?

See 8).

7)\* Prove 6)a).

See 8).

8)\* Characterize the P-positions and N-positions for  $(n_1, n_2, \dots, n_k)$ -Nim. Prove your answer!

$(n_1, n_2, \dots, n_k)$  is a P-position if and only if  $n_1 \oplus n_2 \oplus \dots \oplus n_k = 0$ .

The result follows from the two statements:

i) If  $(n_1, n_2, \dots, n_k)$  is a P-position then any move results in a N-position.

ii) If  $(n_1, n_2, \dots, n_k)$  is a N-position then it is possible to move to a P-position.

i) is straightforward to prove using the Nim-sums.

For ii), suppose  $n_1 \oplus n_2 \oplus \dots \oplus n_k = x \neq 0$ . Let  $i$  be the leftmost nonzero digit of  $x$  in its binary representation. Thus, for some  $j$ , the  $i$ th digit in the binary representation of  $n_j$  is also nonzero. Check that removing  $n_j - (x \oplus n_j)$  from the  $j$ th pile works.

$\infty$ ) Solve the following games (i.e. try to characterize the P-positions and N-positions and try to prove your answer. However, this may not be possible!).

a)  $l$ -Subtraction Game: Similar to  $(n_1, n_2, \dots, n_k)$ -Nim, except players can remove at most  $l$  stones.

$(n_1, n_2, \dots, n_k)$  is a P-position if and only if  $n_1 \pmod{l+1} \oplus n_2 \pmod{l+1} \oplus \dots \oplus n_k \pmod{l+1} = 0$ .

b) Greedy Nim: Similar to  $(n_1, n_2, \dots, n_k)$ -Nim, except players can only remove stones from the largest pile.

A position is a P-position if and only if there are an even number of largest piles.

c) Multiple Pile Nim: Similar to  $(n_1, n_2, \dots, n_k)$ -Nim, except players can also remove the same number of stones from both piles.

Some small p-Positions when  $k = 2$  and  $n_1 \geq n_2$  are  $(0, 0), (2, 1), (5, 3), (7, 4), (10, 6)$ .

Some problems are taken from:

- D. Fomin, S. Genkin, I. Itenberg “Mathematical Circles (Russian Experience)”