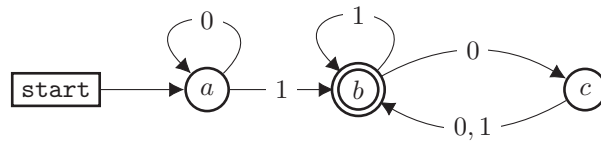

Finite Automata

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Part 1: DFAs

This week, we will study computational devices called *deterministic finite automata*. A DFA has a simple job: it will either “accept” or “reject” a string of letters.

Consider the automaton A shown below:



A takes strings of letters in the alphabet $\{0, 1\}$ and reads them left to right, one letter at a time. Starting in the state a , the automaton A will move between states along the edge marked by each letter.

Note that node b has a “double edge” in the diagram above. This means that the state b is *accepting*. Any string that makes A end in state b is *accepted*. Similarly, strings that end in states a or c are *rejected*.

For example, consider the string 1011.

A will go through the states $a - b - c - b - b$ while processing this string.

Problem 1:

Which of the following strings are accepted by A ?

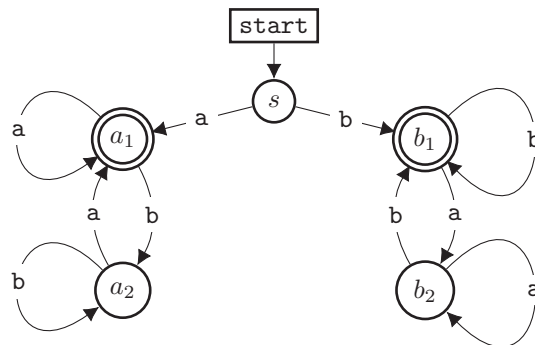
- 1
- 1010
- 1110010
- 1000100

Problem 2:

Describe the general form of a string accepted by A .

Hint: Work backwards from the accepting state, and decide what all the strings must look like at the end in order to be accepted.

Now consider the automaton B , which uses the alphabet $\{a, b\}$. It starts in the state s and has two accepting states a_1 and b_1 .



Problem 3:

Which of the following strings are accepted by B ?

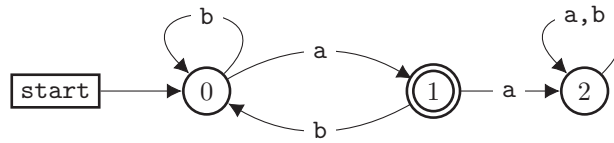
- aa
- abba
- abbba
- baabab

Problem 4:

Describe the strings accepted by B .

Problem 5:

How many strings of length n are accepted by the automaton C ?



Definition 1:

An *alphabet* is a finite set of symbols.

Definition 2:

A *string* over an alphabet Q is a finite sequence of symbols from Q .

We denote the empty string ε .

Q^* is the set of all possible strings over Q .

For example, $\{0, 1\}^*$ is the set $\{\varepsilon, 0, 1, 00, 01, 10, 11, 000, \dots\}$

Note that this set contains the empty string.

Definition 3:

A *language* over an alphabet Q is a subset of Q^* .

For example, the language “strings of length 2” over $\{0, 1\}$ is $\{00, 01, 10, 11\}$

Definition 4:

We say a language L is *recognized* by a DFA if that DFA accepts a string w if and only if $w \in L$.

Problem 6:

Draw DFAs that recognize the following languages. In all parts, the alphabet is $\{0, 1\}$:

- $\{w \mid w \text{ begins with a } 1 \text{ and ends with a } 0\}$
- $\{w \mid w \text{ contains at least three } 1\text{s}\}$
- $\{w \mid w \text{ contains the substring } 0101 \text{ (i.e., } w = x0101y \text{ for some } x \text{ and } y)\}$
- $\{w \mid w \text{ has length at least three and its third symbol is a } 0\}$
- $\{w \mid w \text{ starts with } 0 \text{ and has odd length, or starts with } 1 \text{ and has even length}\}$
- $\{w \mid w \text{ doesn't contain the substring } 110\}$

Problem 7:

Draw a DFA over an alphabet $\{a, b, @, .\}$ recognizing the language of strings of the form $\text{user@website.domain}$, where user , website and domain are nonempty strings over $\{a, b\}$ and domain has length 2 or 3.

Problem 8:

Construct a DFA to check whether a number written in binary is divisible by 2.

Problem 9:

Draw a state diagram for a DFA over an alphabet of your choice that accepts exactly $f(n)$ strings of length n if

- $f(n) = n$
- $f(n) = n + 1$
- $f(n) = 3^n$
- $f(n) = n^2$
- $f(n)$ is a Tribonacci number.

Tribonacci numbers are defined by the sequence $f(0) = 0$, $f(1) = 1$, $f(2) = 1$, and $f(n) = f(n-1) + f(n-2) + f(n-3)$ for $n \geq 3$

Hint: Fibonacci numbers are given by the automaton prohibiting two `a's in a row.

Problem 10:

Draw a DFA recognizing the language of strings over $\{0, 1\}$ in which 0 is the third digit from the end. Prove that any such DFA must have at least 8 states.

Problem 11:

Build a DFA that takes binary strings and accepts it iff it is divisible by 6 when interpreted as a binary whole number (e.g. 10010 is 18 in decimal so it should be accepted. 0 and 0010010 should also be accepted). Read the string from most to least significant digit.